

Assignment 9

Topic ‘Mutation Operators’

In this assignment we always consider a bit string x of length n as an array where the first position is $x[0]$ and, consequently, $x[n-1]$ is the last position.

We consider three different mutation operators, standard bit mutation with mutation probability $p_m = 1/n$, asymmetric mutation, and a mutation operator that we call hypermutation and that is defined as follows.

1. Select $p \in \{0, 1, \dots, n-1\}$ uniformly at random.
2. Select $l \in \{0, 1, \dots, n\}$ uniformly at random.
3. For $i := 0$ to $n-1$ do
4. If $i \leq l-1$
Then $y[(p+i) \bmod n] := 1 - x[(p+i) \bmod n]$
Else $y[(p+i) \bmod n] := x[(p+i) \bmod n]$

Let $m_1(x)$ denote the random bit string resulting from applying standard bit mutation with mutation probability $p_m = 1/n$ to $x \in \{0, 1\}^n$. Let $m_2(x)$ denote the random bit string resulting from applying asymmetric mutation to $x \in \{0, 1\}^n$. Let $m_3(x)$ denote the random bit string resulting from applying hypermutation to $x \in \{0, 1\}^n$.

Problem 1 Consider the following four bit strings (all of length n). $z_1 := 0^n$, $z_2 := 10^{n-1}$, $z_3 := 1010^{n-3}$, $z_4 := 1^n$. For all $i \in \{1, 2, 3\}$, compute $\text{Prob}(m_i(z_1) = z_2)$, $\text{Prob}(m_i(z_2) = z_1)$, $\text{Prob}(m_i(z_1) = z_3)$ and $\text{Prob}(m_i(z_1) = z_4)$.

Problem 2 Consider the (1+1) EA in a variant where you replace standard bit mutation by hypermutation. Prove an upper bound on the expected optimisation time on ONEMAX.

Problem 3 Implementation Task

Remember your implementation for Assignment 8. Add the asymmetric (1+1) EA to your implementation. Perform the same experiments for this RSH as you did for random local search and the (1+1) EA. Compare the results. Can you explain the results?